

A Guaranteed Poly-Logarithmic Time Relaxation for the Line Spectral Estimation Problem

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I. BACKGROUND ON LINE SPECTRAL ESTIMATION

The line spectral estimation problem aims to recover the frequencies of a complex time signal x that is assumed to be sparse in the spectral domain from its discrete measurements $y \in \mathbb{C}^n$, uniformly acquired at a sampling frequency $f_S \in \mathbb{R}^+$. More precisely, the time signal x is assumed to follow the s -spikes model given by

$$\forall t \in \mathbb{R}, \quad x(t) = \sum_{r=1}^s \alpha_r e^{i2\pi\xi_r t}, \quad (1)$$

whereby $\Xi = \{\xi_r\}_{1 \leq r \leq s}$ is the ordered set containing the s spectral components generating the signal x , and $\alpha = \{\alpha_r\}_{1 \leq r \leq s}$ the one of their associated complex amplitudes. The particularity of this model stands in the fact that the frequencies Ξ are drawn *continuously* on $[0, f_S)$ and *are not constrained to belong to some finite discrete grid*, as opposed to discretization-based methods to tackle inverse problems.

This problem is ill-posed and there are infinitely many estimators of the spectral distribution \hat{x} of x that are consistent with the measurement vector y . Among all those estimators, the one considered to be optimal in this spikes recovery context is the one returning a consistent spectral distribution \hat{x}_0 of \hat{x} having the sparsest possible spectral support. Equivalently, this estimator can be defined by the output of the minimization program

$$\begin{aligned} \hat{x}_0 &= \arg \min_{\hat{x} \in D_1} \|\hat{x}\|_0 \\ \text{subject to} \quad & y = \mathcal{F}_n(\hat{x}), \end{aligned} \quad (2)$$

where D_1 denotes the space of absolutely integrable spectral distributions. The functions $\|\cdot\|_0$ and $\mathcal{F}_n(\cdot)$ are respectively the support counting pseudo-norm and the inverse discrete time Fourier transform whose expressions are given in Table I.

Program (2) is non-convex and difficult to solve in a direct approach due to the combinatorial nature of “ L_0 ” minimization. A commonly proposed workaround consists in analysing the output of a *convex relaxation* of (2), obtained by swapping the cardinality cost function $\|\cdot\|_0$ into a minimization of the total-variation norm $|\cdot|(\mathbb{T}_{f_S})$ defined in Table I over D_1 . This relaxation was proven to be tight in [1] and robust to noise in [2], provided that a sufficient separability criterion

$$\Delta_{\mathbb{T}_{f_S}}(\Xi) \geq \frac{4f_S}{n-1} \quad (3)$$

is respected, where $\Delta_{\mathbb{T}_{f_S}}(\cdot)$ is the minimal warp around distance on the rescaled elementary torus $\mathbb{T}_{f_S} = [0, f_S)$ between elements of a set defined in Table I. This bound was tightened later on in [3].

More interestingly, it has been shown in [4] that the tightness of the convex approach still holds with high probability when extracting independently at random a small number of observations and discarding the rest of it. The observation vector $y \in \mathbb{C}^m$ resulting from this random process is linked to the spectrum \hat{x} of the probed signal by the linear relation $y = C_{\mathcal{I}} \mathcal{F}_n(\hat{x})$ where $C_{\mathcal{I}} \in \{0, 1\}^{m \times n}$ is a boolean matrix whose rows are equal to $\{e_k^T\}_{k \in \mathcal{I}}$ and where

$\mathcal{I} \subseteq \llbracket 0, n-1 \rrbracket$ is the subset of cardinality m describing the indexes of the retained samples. In addition, it has been shown that the dual Lagrange program takes the form of the semidefinite program

$$\begin{aligned} (c_*, H_*) &= \arg \max_{\substack{c \in \mathbb{C}^m \\ H \in \mathbb{C}^{n \times n}}} \Re(y^T c) \\ \text{subject to} \quad & \begin{bmatrix} H & q \\ q^* & 1 \end{bmatrix} \succeq 0 \\ & \mathcal{T}_n^*(H) = e_0 \\ & q = C_{\mathcal{I}}^* c, \end{aligned} \quad (4)$$

where \mathcal{T}_n^* is the adjoint of the linear operator \mathcal{T}_n and $\mathcal{T}_n(u)$ is the Toeplitz Hermitian matrix whose first row is equal to u for all $u \in \mathbb{C}^n$. Moreover, the polynomial of degree $n-1$ having for coefficients vector $q_* = C_{\mathcal{I}}^* c_*$ locates with high probability the frequencies supporting \hat{x}_0 around the unit circle.

II. MAIN CONTRIBUTION

The semidefinite program (4) remains of dimension n , which can be much greater than the number of observation m . Its output is computable in $\mathcal{O}(n^7)$ operations via the use of interior point solvers, which become intractable when n exceeds a few hundred. Our result complements the tightness guarantees of [4] by showing the existence of a semidefinite program of dimension m recovering the spectral support of \hat{x}_0 with high probability. Moreover since $m > \mathcal{O}(\log^2 n)$ has been guaranteed in [4] to produce a tight estimate of the spectral support, our program is computable in a *poly-logarithmic time* of the variable n . Our results are summarized by the following theorem, and relies on a novel extension of the theory of Gram parametrization of trigonometric polynomials to subspaces of polynomials [5].

Theorem 1. *Let \mathcal{I} be a subset of cardinality m drawn uniformly at random in $\llbracket 0, n-1 \rrbracket$, and let $\mathcal{R}_{\mathcal{I}}$ the linear operator defined by $\mathcal{R}_{\mathcal{I}}(u) = C_{\mathcal{I}} \mathcal{T}_n(u) C_{\mathcal{I}}^*$ for all $u \in \mathbb{C}^n$. Suppose that x follows Model (1) and satisfies Condition (3). Moreover, suppose that the elements of α have phases drawn independently and uniformly at random in $[0, 2\pi)$. Consider any positive number $\delta > 0$. There exists a constant $C > 0$ such that if*

$$m \geq C \max \left\{ \log^2 \frac{n}{\delta}, s \log \frac{s}{\delta} \log \frac{n}{\delta} \right\},$$

then the semidefinite program

$$\begin{aligned} (c_*, S_*) &= \arg \max_{\substack{c \in \mathbb{C}^m \\ S \in \mathbb{C}^{m \times m}}} \Re(y^T c) \\ \text{subject to} \quad & \begin{bmatrix} S & c \\ c^* & 1 \end{bmatrix} \succeq 0 \\ & \mathcal{R}_{\mathcal{I}}^*(S) = e_0 \end{aligned} \quad (5)$$

outputs with probability greater than $1 - \delta$ a vector $c_* \in \mathbb{C}^m$ such that $q_* = C_{\mathcal{I}}^* c_*$ induces a polynomial Q_* of degree $n-1$ locating the support of \hat{x}_0 . Moreover, this program can be solved in $\mathcal{O}(m^3)$ operations via the alternating direction method of multipliers.

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Function	Domain	Expression
\mathcal{F}_n	$D_1 \rightarrow \mathbb{C}^n$	$\forall k \in \llbracket 0, n-1 \rrbracket, \mathcal{F}_n(\hat{x})[k] = \int_{\mathbb{T}_{f_S}} e^{i2\pi \epsilon k} d\hat{x}(\xi)$
$\Delta_{\mathbb{T}_{f_S}}$	$\wp(\mathbb{T}_{f_S}) \rightarrow \mathbb{R}^+$	$\Delta_{\mathbb{T}_{f_S}}(\Omega) = \inf_{(\xi, \xi') \in \Omega^2} \{ \xi - \xi' : \xi \neq \xi'\}$
$\ \cdot\ _0$	$D_1 \rightarrow \mathbb{R}^+ \cup \{+\infty\}$	$\ \hat{x}\ _0 = \text{card} \{\hat{x}(\xi) \neq 0 : \xi \in \mathbb{T}_{f_S}\}$
$ \cdot (\mathbb{T}_{f_S})$	$D_1 \rightarrow \mathbb{R}^+$	$ \hat{x} (\mathbb{T}_{f_S}) = \sup_{h \in \mathcal{C}(\mathbb{T}_{f_S})} \left\{ \Re \left[\int_{\mathbb{T}_{f_S}} \overline{h(\xi)} d\hat{x}(\xi) \right] : \ h\ _\infty \leq 1 \right\}$

Table I
MATHEMATICAL DEFINITIONS

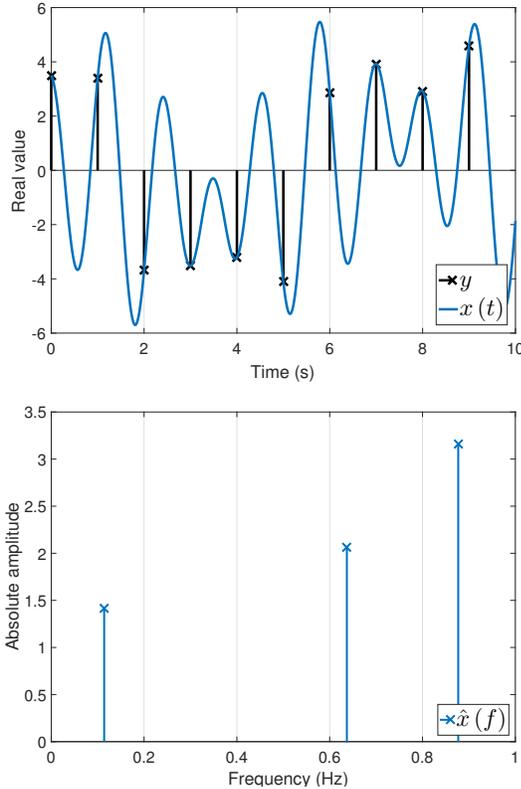


Figure 1. Time and spectral representation of a signal x following the spikes model with three spectral spikes and its measurement vector y when taking $n = 10$ observations at a frequency $f_S = 1\text{Hz}$.

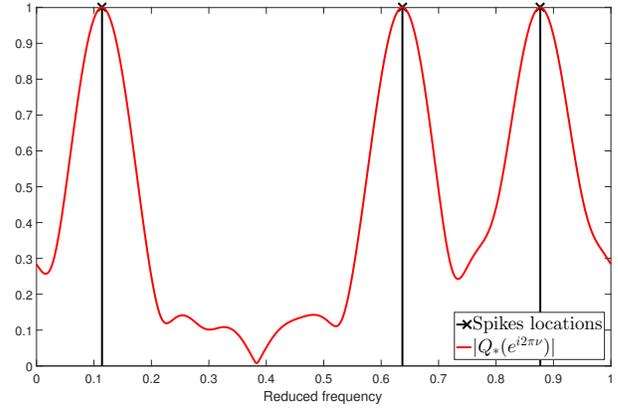


Figure 2. The optimal dual polynomial Q_* obtained by solving Program (5) when retaining entries of y with indexes in the set $\mathcal{I} = \{0, 3, 4, 6, 8, 9\}$ of cardinality $m = 6$. $Q_*(e^{i2\pi\nu})$ locates the frequencies of the signal x by reaching modulus 1 whenever $\nu \in \frac{1}{f_S}\Xi$.