

# A Guaranteed Poly-Logarithmic Time Relaxation for the Line Spectral Estimation Problem

Maxime Ferreira Da Costa and Wei Dai

Department of Electrical and Electronic Engineering, Imperial College London, United Kingdom

Email: {maxime.ferreira, wei.dai1}@imperial.ac.uk

## I. BACKGROUND ON LINE SPECTRAL ESTIMATION

The line spectral estimation problem aims to recover the frequencies of a complex time signal  $x$  that is assumed to be sparse in the spectral domain from its discrete measurements  $y \in \mathbb{C}^n$ , uniformly acquired at a sampling frequency  $f_S \in \mathbb{R}^+$ . More precisely, the time signal  $x$  is assumed to follow the  $s$ -spikes model given by

$$\forall t \in \mathbb{R}, \quad x(t) = \sum_{r=1}^s \alpha_r e^{i2\pi\xi_r t}, \quad (1)$$

whereby  $\Xi = \{\xi_r\}_{1 \leq r \leq s}$  is the ordered set containing the  $s$  spectral components generating the signal  $x$ , and  $\alpha = \{\alpha_r\}_{1 \leq r \leq s}$  the one of their associated complex amplitudes. The particularity of this model stands in the fact that the frequencies  $\Xi$  are drawn *continuously* on  $[0, f_S)$  and *are not constrained to belong to some finite discrete grid*, as opposed to discretization-based methods to tackle inverse problems.

This problem is ill-posed and there are infinitely many estimators of the spectral distribution  $\hat{x}$  of  $x$  that are consistent with the measurement vector  $y$ . Among all those estimators, the one considered to be optimal in this spikes recovery context is the one returning a consistent spectral distribution  $\hat{x}_0$  of  $\hat{x}$  having the sparsest possible spectral support. Equivalently, this estimator can be defined by the output of the minimization program

$$\begin{aligned} \hat{x}_0 &= \arg \min_{\hat{x} \in D_1} \|\hat{x}\|_0 \\ \text{subject to} \quad & y = \mathcal{F}_n(\hat{x}), \end{aligned} \quad (2)$$

where  $D_1$  denotes the space of absolutely integrable spectral distributions. The functions  $\|\cdot\|_0$  and  $\mathcal{F}_n(\cdot)$  are respectively the support counting pseudo-norm and the inverse discrete time Fourier transform whose expressions are given in Table I.

Program (2) is non-convex and difficult to solve in a direct approach due to the combinatorial nature of “ $L_0$ ” minimization. A commonly proposed workaround consists in analysing the output of a *convex relaxation* of (2), obtained by swapping the cardinality cost function  $\|\cdot\|_0$  into a minimization of the total-variation norm  $|\cdot|(\mathbb{T}_{f_S})$  defined in Table I over  $D_1$ . This relaxation was proven to be tight in [1] and robust to noise in [2], provided that a sufficient separability criterion

$$\Delta_{\mathbb{T}_{f_S}}(\Xi) \geq \frac{4f_S}{n-1} \quad (3)$$

is respected, where  $\Delta_{\mathbb{T}_{f_S}}(\cdot)$  is the minimal warp around distance on the rescaled elementary torus  $\mathbb{T}_{f_S} = [0, f_S)$  between elements of a set defined in Table I. This bound was tightened later on in [3].

More interestingly, it has been shown in [4] that the tightness of the convex approach still holds with high probability when extracting independently at random a small number of observations and discarding the rest of it. The observation vector  $y \in \mathbb{C}^m$  resulting from this random process is linked to the spectrum  $\hat{x}$  of the probed signal by the linear relation  $y = C_{\mathcal{I}} \mathcal{F}_n(\hat{x})$  where  $C_{\mathcal{I}} \in \{0, 1\}^{m \times n}$  is a boolean matrix whose rows are equal to  $\{e_k^T\}_{k \in \mathcal{I}}$  and where

$\mathcal{I} \subseteq \llbracket 0, n-1 \rrbracket$  is the subset of cardinality  $m$  describing the indexes of the retained samples. In addition, it has been shown that the dual Lagrange program takes the form of the semidefinite program

$$\begin{aligned} (c_*, H_*) &= \arg \max_{\substack{c \in \mathbb{C}^m \\ H \in \mathbb{C}^{n \times n}}} \Re(y^T c) \\ \text{subject to} \quad & \begin{bmatrix} H & q \\ q^* & 1 \end{bmatrix} \succeq 0 \\ & \mathcal{T}_n^*(H) = e_0 \\ & q = C_{\mathcal{I}}^* c, \end{aligned} \quad (4)$$

where  $\mathcal{T}_n^*$  is the adjoint of the linear operator  $\mathcal{T}_n$  and  $\mathcal{T}_n(u)$  is the Toeplitz Hermitian matrix whose first row is equal to  $u$  for all  $u \in \mathbb{C}^n$ . Moreover, the polynomial of degree  $n-1$  having for coefficients vector  $q_* = C_{\mathcal{I}}^* c_*$  locates with high probability the frequencies supporting  $\hat{x}_0$  around the unit circle.

## II. MAIN CONTRIBUTION

The semidefinite program (4) remains of dimension  $n$ , which can be much greater than the number of observation  $m$ . Its output is computable in  $\mathcal{O}(n^7)$  operations via the use of interior point solvers, which become intractable when  $n$  exceeds a few hundred. Our result complements the tightness guarantees of [4] by showing the existence of a semidefinite program of dimension  $m$  recovering the spectral support of  $\hat{x}_0$  with high probability. Moreover since  $m > \mathcal{O}(\log^2 n)$  has been guaranteed in [4] to produce a tight estimate of the spectral support, our program is computable in a *poly-logarithmic time* of the variable  $n$ . Our results are summarized by the following theorem, and relies on a novel extension of the theory of Gram parametrization of trigonometric polynomials to subspaces of polynomials [5].

**Theorem 1.** *Let  $\mathcal{I}$  be a subset of cardinality  $m$  drawn uniformly at random in  $\llbracket 0, n-1 \rrbracket$ , and let  $\mathcal{R}_{\mathcal{I}}$  the linear operator defined by  $\mathcal{R}_{\mathcal{I}}(u) = C_{\mathcal{I}} \mathcal{T}_n(u) C_{\mathcal{I}}^*$  for all  $u \in \mathbb{C}^n$ . Suppose that  $x$  follows Model (1) and satisfies Condition (3). Moreover, suppose that the elements of  $\alpha$  have phases drawn independently and uniformly at random in  $[0, 2\pi)$ . Consider any positive number  $\delta > 0$ . There exists a constant  $C > 0$  such that if*

$$m \geq C \max \left\{ \log^2 \frac{n}{\delta}, s \log \frac{s}{\delta} \log \frac{n}{\delta} \right\},$$

then the semidefinite program

$$\begin{aligned} (c_*, S_*) &= \arg \max_{\substack{c \in \mathbb{C}^m \\ S \in \mathbb{C}^{m \times m}}} \Re(y^T c) \\ \text{subject to} \quad & \begin{bmatrix} S & c \\ c^* & 1 \end{bmatrix} \succeq 0 \\ & \mathcal{R}_{\mathcal{I}}^*(S) = e_0 \end{aligned} \quad (5)$$

outputs with probability greater than  $1 - \delta$  a vector  $c_* \in \mathbb{C}^m$  such that  $q_* = C_{\mathcal{I}}^* c_*$  induces a polynomial  $Q_*$  of degree  $n-1$  locating the support of  $\hat{x}_0$ . Moreover, this program can be solved in  $\mathcal{O}(m^3)$  operations via the alternating direction method of multipliers.

REFERENCES

- [1] E. J. Candès and C. Fernandez-Granda, "Towards a mathematical theory of super-resolution," *Communications on Pure and Applied Mathematics*, vol. 67, no. 6, pp. 906–956, 2014.
- [2] —, "Super-resolution from noisy data," *Journal of Fourier Analysis and Applications*, vol. 19, no. 6, pp. 1229–1254, 2013.
- [3] C. Fernandez-Granda, "Super-resolution of point sources via convex programming," *Information and Inference*, vol. 5, no. 3, pp. 251–303, Sep. 2016.
- [4] G. Tang, B. N. Bhaskar, P. Shah, and B. Recht, "Compressed sensing off the grid," *IEEE Transactions on Information Theory*, vol. 59, no. 11, pp. 7465–7490, Nov. 2013.
- [5] M. Ferreira Da Costa and W. Dai, "Exact Dimensionality Reduction for Partial Line Spectra Estimation Problems," *arXiv:1609.03142*.
- [6] G. Tang, B. N. Bhaskar, and B. Recht, "Near minimax line spectral estimation," *IEEE Transactions on Information Theory*, vol. 61, no. 1, pp. 499–512, Jan. 2015.
- [7] B. Dumitrescu, *Positive Trigonometric Polynomials and Signal Processing Applications*. Springer, 2010.

Function	Domain	Expression
$\mathcal{F}_n$	$D_1 \rightarrow \mathbb{C}^n$	$\forall k \in \llbracket 0, n-1 \rrbracket, \mathcal{F}_n(\hat{x})[k] = \int_{\mathbb{T}_{f_S}} e^{i2\pi \varepsilon k} d\hat{x}(\varepsilon)$
$\Delta_{\mathbb{T}_{f_S}}$	$\wp(\mathbb{T}_{f_S}) \rightarrow \mathbb{R}^+$	$\Delta_{\mathbb{T}_{f_S}}(\Omega) = \inf_{(\xi, \xi') \in \Omega^2} \{ \xi - \xi'  : \xi \neq \xi'\}$
$\ \cdot\ _0$	$D_1 \rightarrow \mathbb{R}^+ \cup \{+\infty\}$	$\ \hat{x}\ _0 = \text{card} \{\hat{x}(\xi) \neq 0 : \xi \in \mathbb{T}_{f_S}\}$
$ \cdot (\mathbb{T}_{f_S})$	$D_1 \rightarrow \mathbb{R}^+$	$ \hat{x} (\mathbb{T}_{f_S}) = \sup_{h \in \mathcal{C}(\mathbb{T}_{f_S})} \left\{ \Re \left[ \int_{\mathbb{T}_{f_S}} \overline{h(\varepsilon)} d\hat{x}(\varepsilon) \right] : \ h\ _\infty \leq 1 \right\}$

Table I  
MATHEMATICAL DEFINITIONS

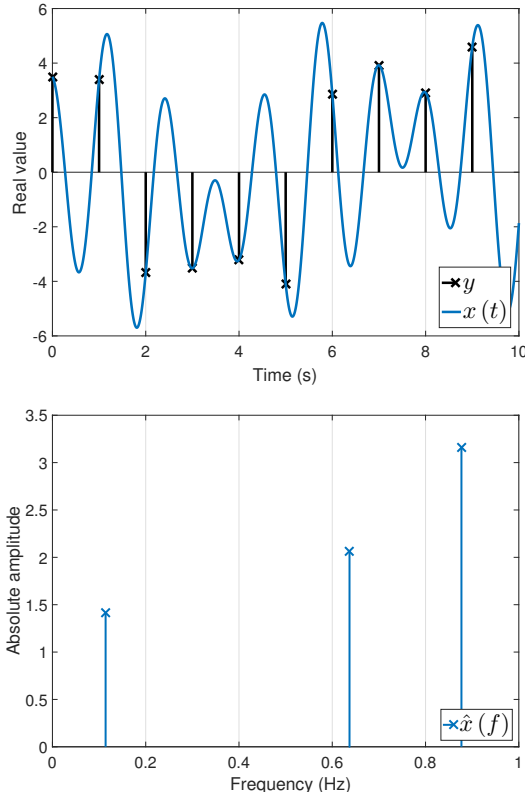


Figure 1. Time and spectral representation of a signal  $x$  following the spikes model with three spectral spikes and its measurement vector  $y$  when taking  $n = 10$  observations at a frequency  $f_S = 1\text{Hz}$ .

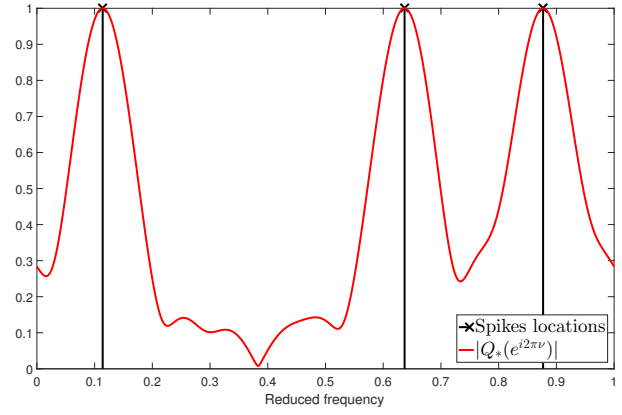


Figure 2. The optimal dual polynomial  $Q_*$  obtained by solving Program (5) when retaining entries of  $y$  with indexes in the set  $\mathcal{I} = \{0, 3, 4, 6, 8, 9\}$  of cardinality  $m = 6$ .  $Q_*(e^{i2\pi\nu})$  locates the frequencies of the signal  $x$  by reaching modulus 1 whenever  $\nu \in \frac{1}{f_S}\Xi$ .